

# The Transport Layer: Review

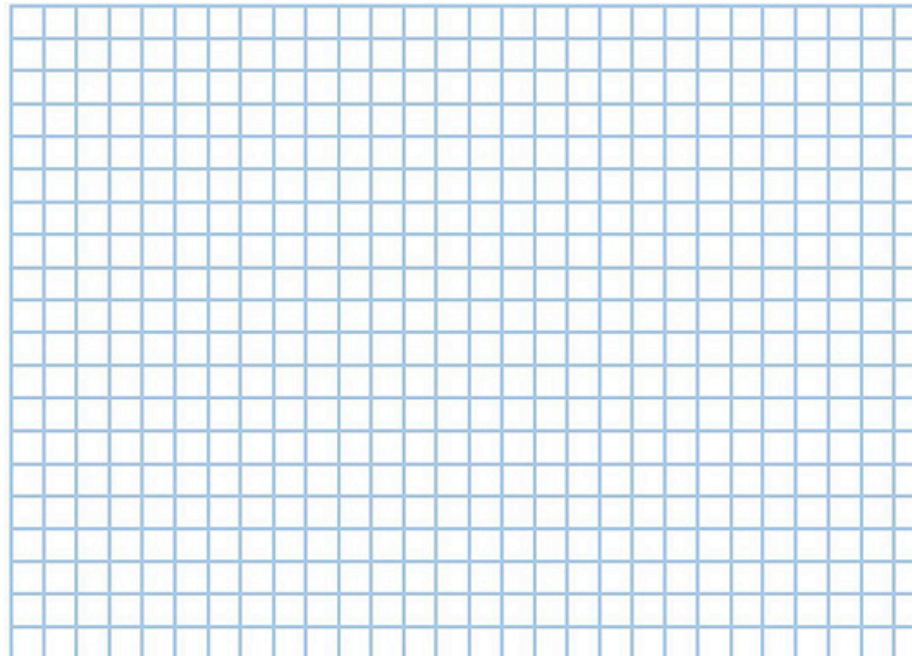
Smith College, CSC 249  
February 22, 2018

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## TCP Congestion Control

1. How does a sender sense congestion?
  - ❖
2. How does a sender determine its sending rate?
  - ❖
  - ❖
3. **What algorithm** is used to change the send-rate?
  - ❖ Many phases and alternatives...

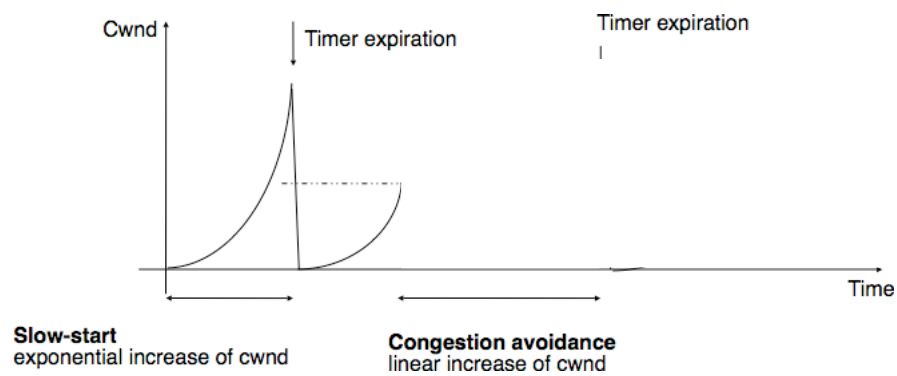
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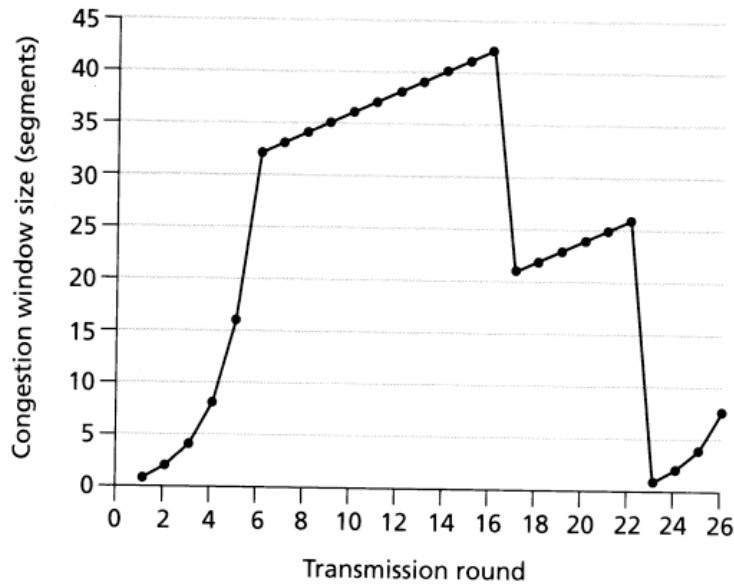
## Reaction to Loss Events

- Exponential increase switches to linear increase when  $\text{CongWin}$  gets to the 'threshold' value (size)



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Identify everything on this graph



## TCP Congestion Control Algorithm

### Increase Sending Rate Phase Options:

1. When CongWin is below Threshold, sender in **slow-start** phase, window grows exponentially.
2. When CongWin is above Threshold, sender is in **congestion-avoidance** phase, window grows linearly.

### Decrease Sending Rate Phase Options:

1. When a **triple duplicate ACK** occurs, Threshold set to CongWin/2 and CongWin set to Threshold.
2. When **timeout** occurs, Threshold set to CongWin/2 and CongWin is set to 1 MSS.

## TCP Congestion Control Algorithm

### Three major phases / mechanisms:

- 1) Slow start - at 1 max segment size
  - But increase \_\_\_\_\_
- 2) Congestion Avoidance phase
  - AIMD = additive incr, multiplicative decr
  - Using `cwnd` and `ssthresh`
- 3) Fast Recovery
  - Increase of `cwnd` each round trip time
  - Slow start: \_\_\_\_\_
  - Congestion avoidance: \_\_\_\_\_

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## Summary TCP reaction to loss

- Loss indicated by timeout
  - `cwnd` set to 1 MSS
  - `ssthresh` set to `cwnd`/2
  - Window (`cwnd`) grows exponentially (slow start) to the threshold, then grows linearly
- Loss indicated by 3 duplicate ACKs
  - Network capable of delivering some segments, so...
  - `cwnd` is cut in half (=ssthresh)
  - Window grows linearly

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## Transport Layer Review

□ The transport layer services are:

- ❖
- ❖
- ❖
- ❖
- ❖
- ❖
- ❖
- ❖

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## Transport Layer Review

□ The transport layer ***does not*** provide:

- ❖
- ❖
- ❖
- ❖
- ❖
- ❖
- ❖
- ❖

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## Transport Layer Review

- ❑ Compare TCP and UDP (pros and cons?)



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## Transport Layer Review

- ❑ TCP Connection Management includes



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## Transport Layer Review

### □ Elements of TCP reliability:

- ❖
- ❖
- ❖
- ❖
- ❖
- ❖
- ❖
- ❖

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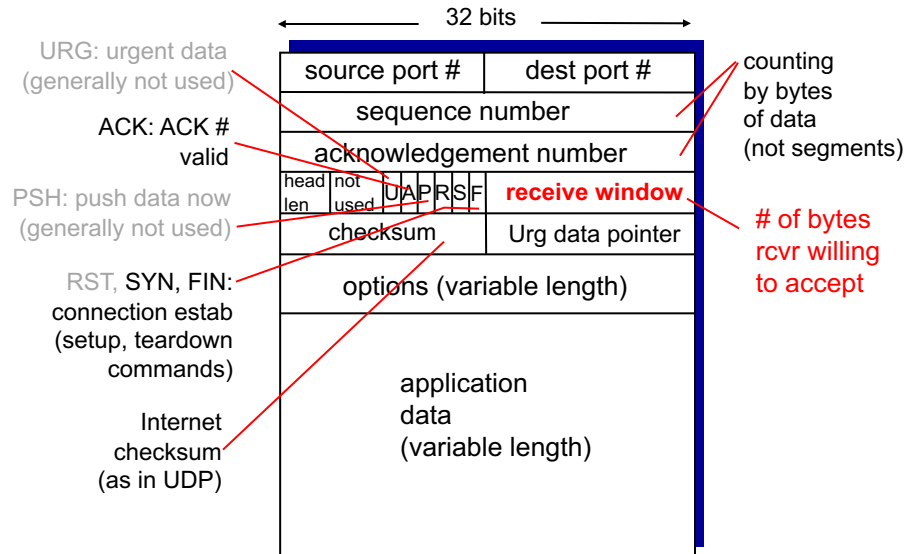
## Transport Layer Review

### □ Elements of congestion control algorithm

- ❖
- ❖
- ❖
- ❖
- ❖
- ❖
- ❖
- ❖

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## TCP Flow Control



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## TCP flow control (quick & easy)

- **Receiver** “advertises” free buffer space by including **rwnd** value in TCP header of receiver-to-sender segments
  - ❖ **RcvBuffer** size set via socket options (typical default is 4096 bytes)
  - ❖ Many operating systems auto-adjust **RcvBuffer**
- **Sender** limits amount of un-ACKed (“in-flight”) data to receiver’s **rwnd** value
  - ❖ Guarantees receive buffer will not overflow

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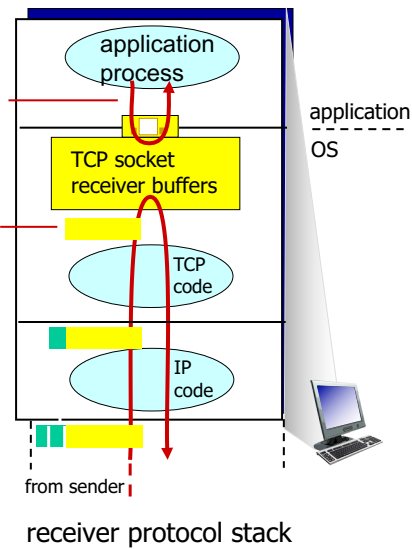
## TCP flow control

The application may remove data from the TCP socket buffers ....

... slower than the TCP receiver is delivering it into the buffers (sender is sending)

*flow control to the rescue!*

receiver controls sender, so sender won't overflow receiver's buffer by transmitting too much, too fast



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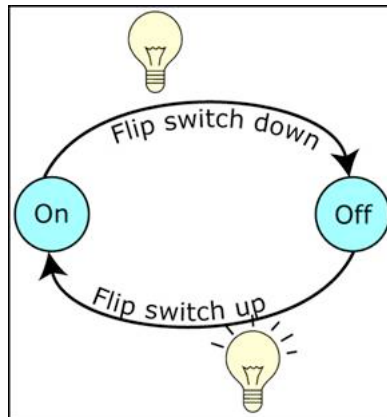
## Transport Layer Review

❑ Other questions?

- ❖
- ❖
- ❖
- ❖
- ❖
- ❖
- ❖
- ❖

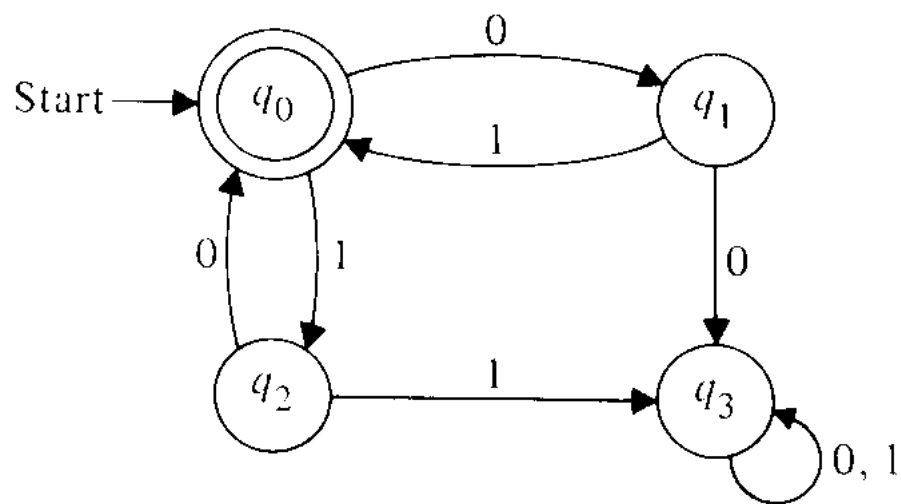
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## A fun tangent... Finite State Machines



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## Finite State Machines



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## TCP sender events:

### (1) data received from application:

1. Create a segment and assign a SEQ number
  - ❖ SEQ # is byte-stream number of first data byte in segment
2. Start timer if it is not already running
  - ❖ Timer is for the oldest un-acked segment
  - ❖ Expiration interval: `TimeoutInterval`

### (2) timeout:

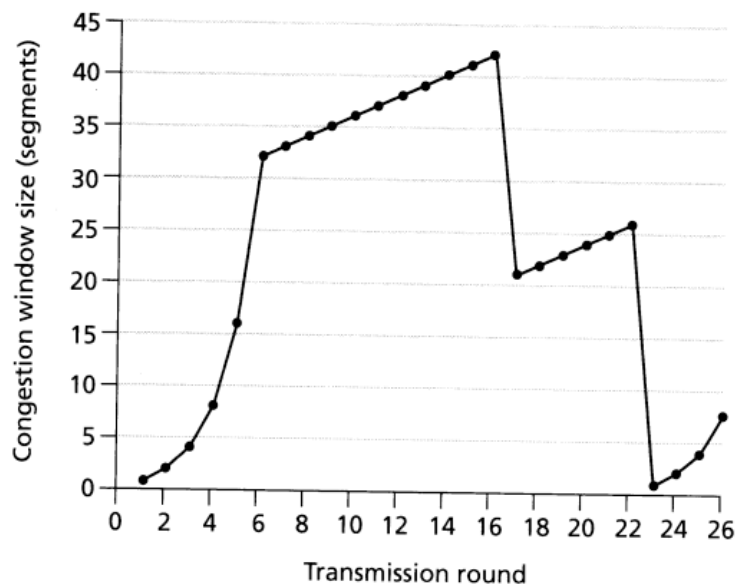
1. Retransmit segment that caused the timeout
2. Restart the timer

### (3) ACK received:

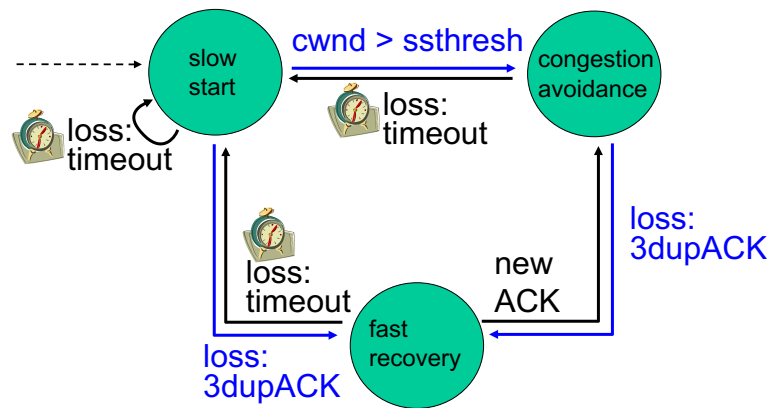
- ❑ For previously unacked segments
  1. update what is known to be acked
  2. start timer if there are outstanding segments

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## TCP Congestion Control



## TCP Congestion Control: FSM



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## Transport Layer Summary

- ❑ TCP and UDP Services
- ❑ Encapsulation (create and attach header)
- ❑ Multiplexing and demultiplexing
- ❑ Checksum
- ❑ Connection management
- ❑ Reliable transport service
- ❑ Congestion control
- ❑ Detect loss and retransmit
  - ❖ Detect out-of-order and reorder
- ❑ Flow Control

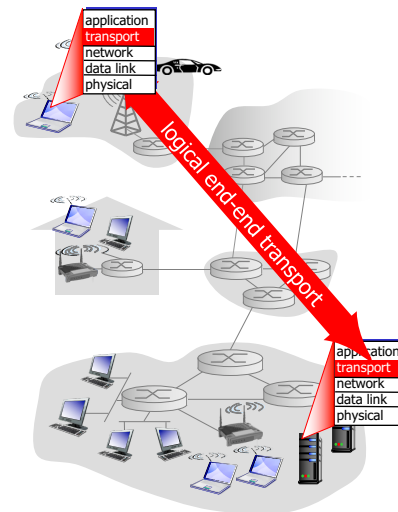
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## Transport services and protocols

- Provide *logical communication, a virtual connection*

...between application **processes** running on different hosts

*This is not* a physical path including routers



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## Introduction to the Network Layer

- Desired network layer services...
  - Actual network layer services
- Implemented in hosts and routers
- Two main network layer functions
- Three main network layer protocols

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## Network Layer Services of IP?

- ❑ Guaranteed delivery?
- ❑ Guaranteed minimum delay?
- ❑ In-order datagram delivery?
- ❑ Guaranteed minimum bandwidth to flow?
- ❑ Restrictions on changes in inter-packet spacing?
- ❑ IP Provides? → “Best-effort service”

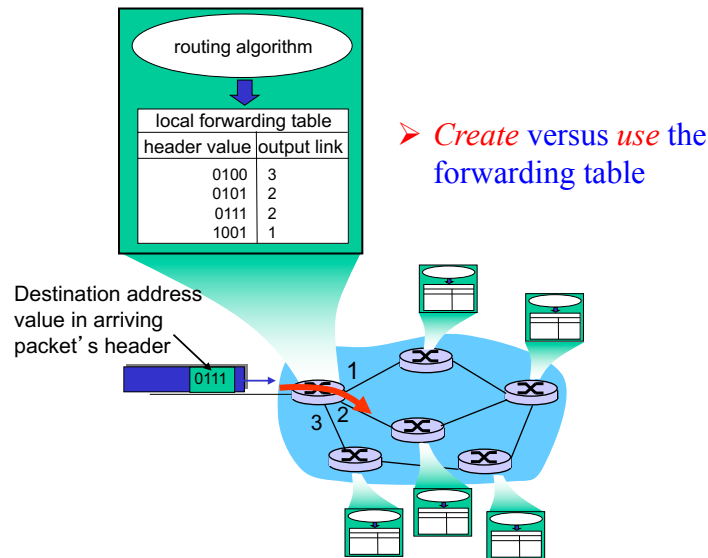
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## Key Network-Layer Functions

1. *routing*: determine route taken by packets from source to destination
  - ❖ *Network-wide routing algorithms*
2. *forwarding*: move packets from router's input link to appropriate output link
  - ❖ *Internal to a single router*

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## Network Layer: Routing and Forwarding



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## Network Layer, Chapter 4

- ❑ Router 'switching fabric'
  - ❖ Hardware / electrical pathways within a router
- ❑ Forwarding - use the forwarding table to transmit, or forward, each packet to the correct output link, based on the destination IP address
- ❑ Routing - Create the forwarding tables
  - ❖ Decentralized vs. Centralized algorithms
  - ❖ Within an ISP vs. between ISPs
- ❑ Software Defined Networks, SDN

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