

## Overview: TCP Basics

- □ Recap: using SEQ and ACK numbers
  - \* SEQ random initial number for numbering the bytes in the application message
  - \* "ACKnowledge" <u>next</u> byte expected
  - \* Read chapter to review

#### □ TCP congestion control

\* Read through multiple times!

TCP flow control

# TCP: SEQ and ACK numbers



TCP: Cumulative ACK



3-4

# TCP possible sender events:

#### (1) Data received from application:

- 1. Create a segment and assign a SEQ number
  - SEQ # is byte-stream number of first data byte in segment
- 2. Start timer if it is not already running
  - Timer is for the oldest un-acked segment
  - Expiration interval: TimeOutInterval
- (2) <u>Timeout (ACK not received)</u>:
  - 1. Retransmit segment that caused the timeout
  - 2. Restart the timer

#### (3) ACK received for previously unacked segments

- 1. Update what is known to be acked
- 2. Start timer if there are outstanding segments

5

6



#### TCP: retransmission scenarios



<u>New Today:</u> Principles of Congestion Control

- Packet loss is caused by overflowing router buffers
- Retransmission treats the symptom
- Congestion control treats the cause
- □ What are costs of congestion?
  - \*
  - \*
  - \*

# TCP Congestion Control

- Three questions
  - 1. How does a sender sense congestion?
  - 2. How does a sender limit its sending rate?
  - 3. What algorithm is used to change the send-rate?

# TCP Congestion Control: details

How does sender perceive congestion?

□ A loss event is ?

A timeout or

3 duplicate ACKs

#### How does sender limit its send rate?

TCP sender reduces the send rate via changing a variable value, the "CongWin," after a loss event

LastByteSent-LastByteAcked < CongWin</pre>

# **TCP Congestion Control Window**

- CongWin is dynamic, function of perceived network congestion
- Sender limits transmission: LastByteSent-LastByteAcked < CongWin</p>



□ Changing CongWin changes the SendRate

13

usable, not yet sent

not usable

# TCP Congestion Control Algorithm



# TCP Congestion Control Algorithm

Three major mechanisms:

- 1) Slow start
- 2) Congestion Avoidancce
  - AIMD = additive increase, multiplicative decrease
- 3) Fast Recovery:
  - Reaction to timeout events versus
    3 duplicate ACKs

Are used to adjust ... ?

Adjust CongWin

TCP Slow Start

- 1) When connection begins, CongWin = 1 MSS
  - Available bandwidth probably much greater
  - \* Desirable to quickly ramp up to respectable rate
- Increase rate exponentially fast until first loss event
  - \* Grow window 1MSS for each ACK received
- Summary: initial rate is slow but ramps up exponentially fast

### TCP: detecting, reacting to loss

- Loss indicated by timeout
  - cwnd set to 1 MSS;
  - Window (cwnd) grows exponentially (slow start) to the threshold, then grows linearly
- Loss indicated by 3 duplicate ACKs
  - Duplicate ACKs indicate network capable of delivering some segments
  - cwnd is cut in half window then grows linearly

#### TCP Slow Start with Exp. Increase

- When connection begins, increase rate exponentially until first loss event:
  - ✤ Double Cong₩in every RTT
  - Increment CongWin for every ACK received



## TCP: Congestion algorithm switching from slow start to CA

Implementation:

- variable ssthresh
- on loss event, ssthresh is set to 1/2 of cwnd just before loss event



19

#### Reaction to Loss Events

Exponential increase switches to linear increase when CongWin gets to the 'threshold' value (size)



### Identify everything on this graph



### Summary: TCP Congestion Control

Increase Sending Rate Phase Options:

- 1. When CongWin is below Threshold, sender in slow-start phase, window grows exponentially.
- When CongWin is above Threshold, sender is in congestion-avoidance phase, window grows linearly.

Decrease Sending Rate Phase Options:

- When a triple duplicate ACK occurs, Threshold set to CongWin/2 and CongWin set to Threshold.
- 2. When timeout occurs, Threshold set to CongWin/2 and CongWin is set to 1 MSS.