

The Network Layer: Routing 2: Distance Vector

Smith College, CSC 249
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Overview

- Routing Algorithms
 - ❖ Link-state - From last week
 - ❖ Distance-vector - TODAY

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Overview of Routing so far

- ❖ Routing algorithms
 - ❖ Find the 'best' path through a network
 - ❖ Create forwarding tables
- ❖ Routing occurs between routers (not hosts)
- ❖ Differences between centralized (global) and decentralized algorithms
 - ❖ What are examples of each
 - ❖ Amount of information known initially
 - ❖ How information is shared/spread
 - ❖ Synchronous or asynchronous?

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Algorithm 2: Distance Vector

Rather than using global information, a distance vector algorithm is:

- ❑ **distributed:**
 - ❖ each node communicates only with directly-attached neighbors
- ❑ **iterative:**
 - ❖ continues until no nodes exchange info.
 - ❖ self-terminating: no "signal" to stop
- ❑ **asynchronous:**
 - ❖ nodes need *not* exchange information or iterate in lock step!

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Distance Vector Algorithm

Bellman-Ford Equation

Define

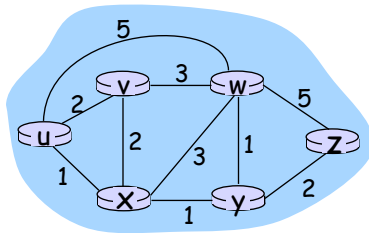
$d_x(y) :=$ cost of least-cost path from x to y

Then

$$d_x(y) = \min_v \{ c(x,v) + d_v(y) \}$$

where min is taken over all neighbors v of x

Bellman-Ford Equation



Clearly, $d_v(z) =$, $d_x(z) =$, $d_w(z) =$

B-F equation says:

$$d_u(z) = \min \{ c(u,v) + d_v(z), \\ c(u,x) + d_x(z), \\ c(u,w) + d_w(z) \}$$

$$= \min \{$$

$$\} =$$

The node that achieves the minimum, is the next hop in the shortest path → forwarding table

Distance Vector Routing Algorithm

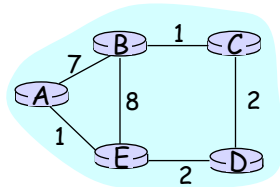
Distance Table data structure

- each node has
 - ❖ A row for each possible destination
 - ❖ A column for each directly-attached neighbor
- example: in node X, for destination Y via neighbor Z:

$$D^X(Y,Z) = \begin{array}{l} \text{distance from X to Y,} \\ \text{via Z as next hop} \\ = c(X,Z) + \min_w \{D^Z(Y,w)\} \end{array}$$

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Distance Table: example with complete information



$$D^E(C,D) = c(E,D) + \min_w \{D^D(C,w)\}$$

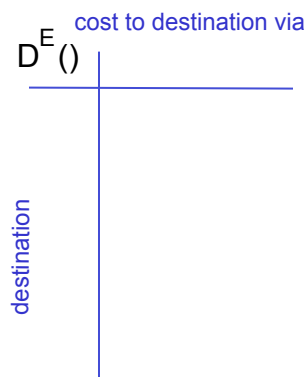
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$$D^E(A,D) = c(E,D) + \min_w \{D^D(A,w)\}$$

$$=$$

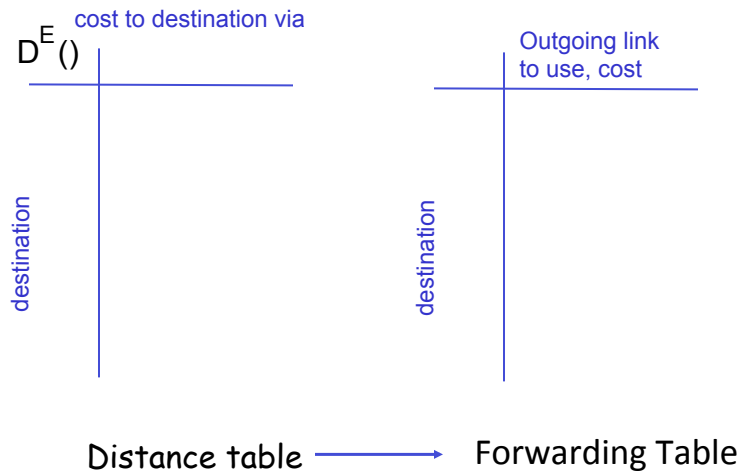
$$D^E(A,B) = c(E,B) + \min_w \{D^B(A,w)\}$$

$$=$$



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Distance table to forwarding table



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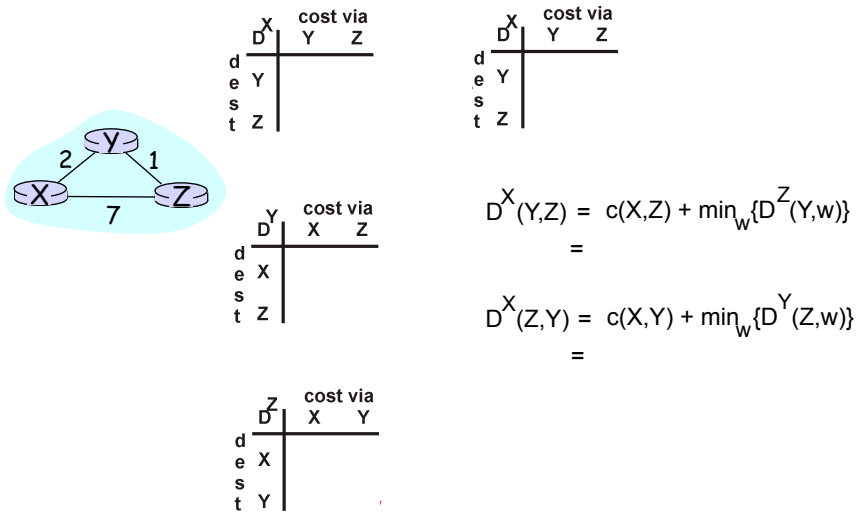
Distance vector algorithm

Asynchronous Iterations:

- Each node begins with $D_x(y)$
 - ❖ An estimate of the cost of the least-cost path from itself to node y , for all nodes in N
- Each node periodically sends its own distance vector estimate to neighbors
 - ❖ → A vector of least costs from itself to all routers
- When a node x receives new DV estimate from neighbor, it updates its own DV using B-F equation, and sends any update to its neighbors
 - $D_x(y) \leftarrow \min_v \{c(x,v) + D_v(y)\}$ for each node $y \in N$
- Under normal conditions, the estimate $D_x(y)$ converges to the actual least cost $d_x(y)$

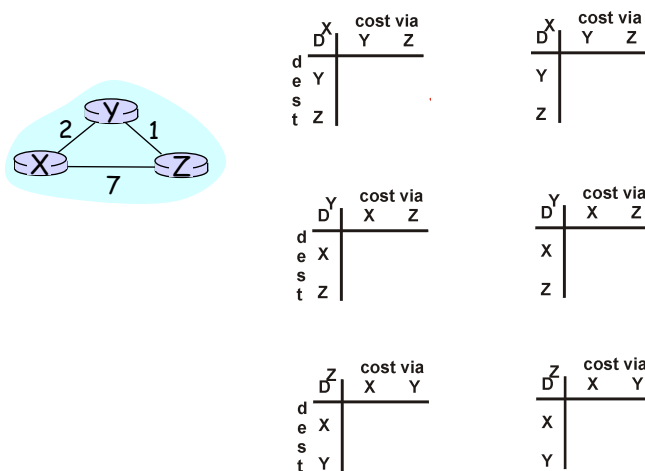
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Distance Vector Algorithm: example for obtaining complete information



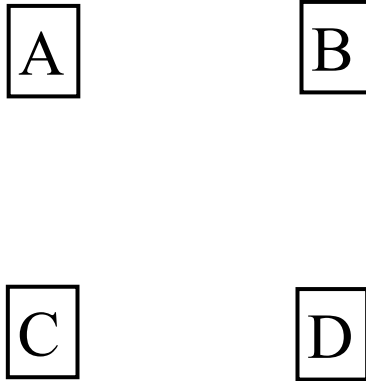
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Distance Vector Algorithm: obtaining info



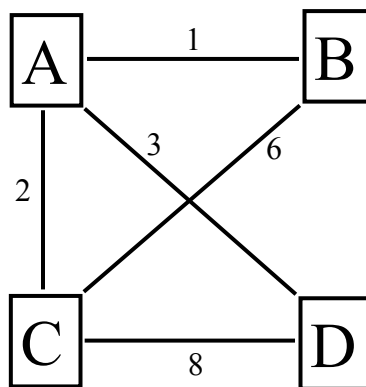
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Distance Vector Routing Activity



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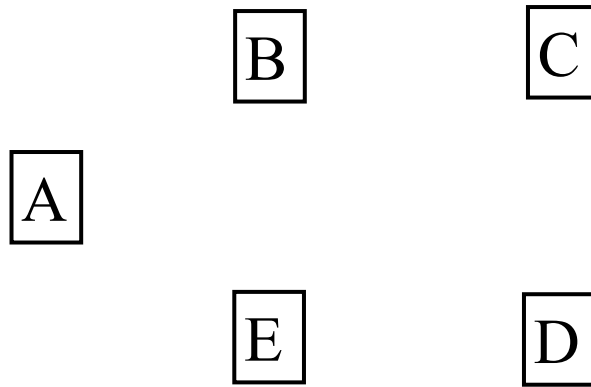
Distance Vector Routing Activity



- Review actual graph – does it match activity results?
- What happens if/when $c(A,D) = 4$ & $c(C,D) = 1$?

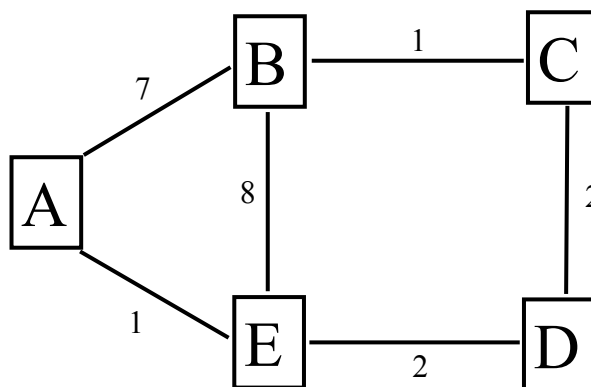
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Distance Vector Routing Activity



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Distance Vector Routing Activity



- Review actual graph – does it match your results?
- What happens if/when $c(A,B) = 2$ and/or if $c(C,D) = 5$?

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Comparison of LS and DV algorithms

- ❑ Information requirements
- ❑ Message complexity
- ❑ Convergence time varies
- ❑ Robustness: what happens if router malfunctions?
- ❑ Oscillations possible?
- ❑ Loops possible?

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Summary

Forwarding:

- ❑ Leads to questions of addressing
 - ❖ Assignment of IP addresses
 - ❖ NAT, IPv6 ...

Routing:

- ❑ Routing objectives
- ❑ Routing notation
- ❑ Routing classification
- ❑ Link state v. Distance Vector
- ❑ Hierarchical structure

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